Needle is an object-oriented functional programming language with a multimethod-based OO system, and a static type system with parameterized types and substantial ML-style type inference.

### The Static Typing Heresy

In exploratory programming:

- Bottom-up programming; design is discovered, not imposed
- Domain knowledge comes from trying to solve the problems
- Subproblems knitted together to build abstractions

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Exploratory programming is easier in a statically typed language with generic functions!

What are generic functions and multimethods?

First, let's look at classes in Needle.

class Thing {} // define a root class

class Rock(Thing) {}
class Paper(Thing) {}
class Scissors(Thing) {}

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```

The curlies are simply where you define members:

```
class Cons[a] (List) {
   head a;
   tail List[a];
}
```

```
generic beats? (Thing, Thing) -> Boolean;
```

```
method beats? (x Rock, y Scissors) { true; }
method beats? (x Paper, y Rock) { true; }
method beats? (x Scissors, y Paper) { true; }
method beats? (x Thing, y Thing) { false; }
```

beats?(rock, rock)  $\Rightarrow$  false

In traditional OO, adding new methods to a class is unmodular even if it's possible.

generic inflammable? Thing -> Boolean;

method inflammable? (x Thing) { false }
method inflammable? (x Paper) { true }

Higher-order functions easily parameterize over behavior, but they don't parameterize over similar data types very well.

In Scheme:

(map function sequence) ;; for lists
(vector-map function sequence) ;; for vectors
(string-map function sequence) ;; for strings

In Needle:

```
generic map c < Sequence . (a -> b, c[a]) -> c[b];
```

- Inspired by Bourdoncle and Merz's ML-sub (1997) and Bonniot (2001)
- Supports parametric types
- Supports type inference

generic map c < Sequence . (a -> b, c[a]) -> c[b];

Type composed of two parts:

- ML-style type scheme
- Type constraints

Type scheme

A type scheme is:

- A nonpolymorphic class Rock, Boolean
- any of a set of type variables a, b, c
- A filled-in polymorphic class List[List[Int]], a  $\rightarrow$  Boolean

Type constraints are conjunctions of subtype relationships; limit which types are permitted to satisfy the type variables in the type scheme.

```
generic negate a < Number . a -> a;
generic map c < Sequence . (a -> b, c[a]) -> c[b];
```

```
fun(seq) { map(negate, seq) }
// has type c < Sequence & a < Number . c[a] -> c[a]
```

- 1. Top-down walk of each top-level expression's AST
- 2. Generate types of subexpressions, combining their constraint sets.
- 3. Verify the constraints are satisfiable
- 4. Simplify the constraints

We do type inference on all code that isn't a generic declaration, or the argument specializer list on a method.

```
let foo = fun (x, f, g, seq) {
    if (x == 3) {
        map(f, seq);
    } else {
        map(fun (x) { g(g(x)) }, seq)
     };
```

has inferred type:

c < Sequence . (Integer, a -> a, a -> a, c[a]) -> c[a] The Needle Programming Language - Neel Krishnaswami - neelk@alum.mit.edu The raw, unsimplified type of foo:

h < Number & (h, h) -> Boolean < (a, Integer) -> g & j < Sequence & (k -> l, j[k]) -> j[l] < (b, d) -> i & n < Sequence & (o -> p, n[o]) -> n[p] < (c, d) -> m & (Boolean, f, f) -> f < (g, i, m) -> e . (a, b, c, d) -> e

After simplification:

c < Sequence . (Integer, a -> a, a -> a, c[a]) -> c[a]

- Generics support separate compilation
- Coverage/completeness tests independent of typechecking
- Everything but generics and method specializers have types inferred

How does the combination of type inference and generics really help?

Best way of composing subproblems is higher-order functions. Static typing helps here, because:

- Easier catch errors when the compiler fails fast
- Easier to discover common patterns when you can see the types
- Generic functions reduce necessary number of parameters

A type is a partial, approximate specification of a function.

Type inference means the compiler generate summaries for you.

## Bottom-up style; design is discovered, not invented

- ML makes it hard to extend datatypes, but easy to write new behaviors
- Java makes it hard to add behaviors, but easy to extend datatypes
- Functions grow behavior; classes grow data
- For exploratory programming you need both

## Future Work

- Improve type simplification
- Add dynamic loading
- Add interfaces

In current Needle, generic printing might have the interface:

generic print a -> String;

```
method print (s String) { s }
method print (b Boolean) { if (b) { "true" } else { "false" } }
```

method print (o a) { raise Error(); }

Throwing an exception hurts safety.

What we want is something like this:

```
interface Print(a) {
    print a -> String;
}
```

generic print Print(a) . a -> String

String implements Print; // interfaces are added \*post-hoc\*
Boolean implements Print;

```
method print (s String) { s }
method print (b String) { if (b) { "true" } else { "false" } }
```

# Interfaces

- Lets you add existing types to new protocols
- Fixes weakness of generic-function style grouping methods.
- Idea stems from Haskell typeclasses.
- Implementation in progress.

- Website at: http://www.nongnu.org/needle
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